Intermediate MPI (Message-Passing Interface)

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Suppose process 0 wants to send a message to process 1. Three possibilities:

 Process 0 can stop and wait until Process 1 is ready to receive the message.

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- Process 0 can stop and wait until Process 1 is ready to receive the message.
- Process 0 can copy the message into a buffer (internal to the library or user-specified) and return from the MPI_Send call.

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It can report failure.

An MPI implementation is allowed to use the first or second interpretation but is not required to use the second one.

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- Use buffered sends. We provide the buffering.
- Use nonblocking communication. This can often give the best performance, especially if we use it to overlap communication and computation.
- Use synchronous sends. MPI provides MPI_Ssend. Send doesn't return until the destination process starts receiving the message. However, this can have performance and scalability issues.

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- Synchronous Mode Send and receive can start before each other but can only complete together.
- Ready Mode Send can only start if matching receive already reached, otherwise error. Use with care.

- Each of the four modes can be applied to both blocking and nonblocking send routines.
- Only the standard mode is available for the blocking and nonblocking receive routines.
- Any type of send routine can be used with any type of receive routine.

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Prototypes.

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 MPI_Bsend allows the user to send messages without worrying about where they are buffered (because the user must have provided buffer space with MPI Buffer attach).

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- MPI_Bsend allows the user to send messages without worrying about where they are buffered (because the user must have provided buffer space with MPI Buffer attach).
- The buffer size given should be the sum of the sizes of all outstanding Bsends that you intend to have, plus MPI_BSEND_OVERHEAD for each Bsend that will be done.

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- The buffer size given should be the sum of the sizes of all outstanding Bsends that you intend to have, plus MPI_BSEND_OVERHEAD for each Bsend that will be done.
- MPI_Buffer_detach returns the buffer address and size so that nested libraries can replace and restore the buffer.

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- Nonblocking send MPI_lsend(...) will return "immediately" even before source location is safe to be altered.
- Nonblocking receive MPI_Irecv(...) will return even there is no message to accept.

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Nonblocking routine formats

MPI_lsend(buf,count,datatype,dest,tag,comm,request)

MPI_Irecv(buf,count,datatype,source,tag,comm, request)

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Completion detected by MPI_Wait() and MPI_Test().

Nonblocking routine formats

MPI lsend(buf,count,datatype,dest,tag,comm,request) MPI lrecv(buf,count,datatype,source,tag,comm, request) Completion detected by MPI Wait() and MPI Test(). MPI Wait(MPI Request *request, MPI Status *status) int MPI Waitall(int count, MPI Request array of requests[], MPI Status array of statuses[]) int MPI Waitany(int count, MPI Request array of requests[], int *index, MPI Status *status) MPI Test(MPI Request *request, int *flag, MPI Status *status)

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```
MPI_Comm_rank(MPI_COMM_WORLD, &myrank); /* find rank */
if (myrank == 0) {
    int x;
    MPI_Isend(&x,1,MPI_INT, 1, msgtag, MPI_COMM_WORLD, req1);
    compute();
    MPI_Wait(req1, status);
} else if (myrank == 1) {
    int x;
    MPI_Recv(&x,1,MPI_INT,0,msgtag, MPI_COMM_WORLD, status);
}
```

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Sending/Receiving structures (Part 1)

▶ We can send a structure by packing it as an array of bytes:

```
struct test {
    int n;
    double x[3], y[3];
};
if (pid == source) {
    struct test test1;
    MPI_Send(&test1, sizeof(struct test), MPI_BYTE, destination,
              tag, MPI_COMM_WORLD);
} else (pid == destination) {
    struct test test2;
    MPI_Recv(&test2, sizeof(struct test), MPI_BYTE, source, tag,
              MPI_COMM_WORLD, status);
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```
However, this relies on the layout of the structure being the same on all nodes.
It also obfuscates the code and introduces platform dependency so it is not a
recommended practice for MPI programs.
```

Sending/Receiving structures (Part 2)

▶ We can send a structure by creating a custom MPI data type for the structure

```
struct test {
       int n;
       double x[3], y[3];
};
const int nitems = 3;
int blocklengths [3] = \{1, 3, 3\}; //lengths of i, x and y as #items
MPI_Datatype types[3] = {MPI_INT, MPI_DOUBLE, MPI_DOUBLE};
MPI_Aint offsets[3];
MPI_Datatype mpi_test_type;
offsets[0] = offsetof(struct test, n);
offsets[1] = offsetof(struct test, x);
offsets[2] = offsetof(struct test, y);
MPI_Type_create_struct(nitems, blocklengths, offsets, types,
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See example lab/MPI/send-struct/ for a working example. This is the recommended way of sending a structure in MPI.

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- See example lab/MPI/send-struct/ for a working example. This is the recommended way of sending a structure in MPI.
- Note that there is no way to send a structure that has variable length (because of pointers stored in it) in one message. We have to use two messages.