A Library of Parallel Algorithms for Generating Combinatorial Objects

Amit Jain Dept. of Math. and Computer Sc. Boise State University Boise, Idaho Elizabeth Elzinga Dept. of Computer Science Grand Valley State University Allendale, Michigan

Background

Sequential algorithms.

- Combinatorial Algorithms by Albert Nijenhuis and Herbert S. Wilf. (2nd ed., 1978)
- The packages Combinatorica for Mathematica and Combinat for Maple.
- The Stony Brook Algorithm Repository, by Steven S. Skiena Home page: http://www.cs.sunysb.edu/~algorith

Parallel algorithms have been published for generating some of the combinatorial objects in research papers but there are no available implementations.

Initial Project Goals

- all/random subsets of an n-set.
- all/random k-subsets of an n-subset.
- all/random permutations of n objects.
- all/random partitions of an *n*-subset.

Underlying Hardware Technology

- Beowulf cluster of 8 dual-Pentium PCs with two 100 Mbs Fast Ethernet networks. The cluster is isolated from the rest of the department via a switch. Check out the book *How to Build a Beowulf* that just came out from MIT press.
- Beowulf home page: http://www.beowulf.org

Underlying Software Technology

- PVM (Parallel Virtual Machine) software. This software is based on message passing between processes which can be on any heterogenous collection of machines. The software runs on almost all kinds of machines. Hence any program written with PVM can be used by almost any combination of hardware platform and operating systems.
- PVM home page: http://www.epm.ornl.gov/pvm_home.html

Project Accomplishments

- Generate all subsets of an n-set. Three different algorithms have been implemented. These are:
 - A *direct* algorithm.
 - A *lexicographic* algorithm.
 - A binary reflected gray code based algorithm.

The sequential generation algorithms were taken from **Wilf**. The algorithms for ranking and unranking (needed for the parallel generation) were designed by us.

- Generate k-subsets of an n-set. Two different approaches were implemented.
 - A *lexicographic* algorithm.
 - A binary reflected gray code based algorithm.

The sequential generation algorithms were taken from **Wilf**. The algorithms for ranking and unranking (needed for the parallel generation) were designed by us for the gray code based algorithm. The ranking and unranking algorithms for the lexicographic generation were based on **Akl**. We did discover an error in the published algorithm that was communicated to the author.

- Generate all permutations of $\{1, 2, ..., n\}$. Two different parallel algorithms were implemented.
 - A *lexicographic* algorithm.
 - A binary reflected gray code based algorithm.

The sequential gray-code based algorithm was taken from **Wilf**. The algorithms for ranking and unranking (needed for the parallel generation) were designed by us. The ranking and unranking algorithms for the lexicographic generation were based on the algorithms published in **Akl**.

- Wilf Combinatorial Algorithms by Albert Nijenhuis and Herbert S. Wilf. (Academic Press, 2nd ed., 1978)
- **Akl** The Design and Analysis of Parallel Algorithms by S. G. Akl (Prentice Hall, 1989). Also the papers referred to in the book in Chapter 6.

	size	sequential time	parallel time	speedup
	$(2^{size}subsets)$	(seconds)	(seconds)	(with 16 CPUs)
direct:				
	22	1.14	0.20	5.7
	24	4.57	0.55	8.3
	26	18.25	1.93	9.5
	28	72.85	7.48	9.7
	30	291.26	30.11	9.7
	32	1165.89	113.79	10.2
lexicographic:				
	22	1.35	0.19	7.1
	24	4.77	0.53	9.0
	26	16.98	1.85	9.2
	28	67.90	7.16	9.5
	30	271.62	29.44	9.2
	32	1088.10	108.84	10.0
graycode:				
	22	1.02	0.19	5.4
	24	4.13	0.52	7.9
	26	16.57	1.92	8.6
	28	65.87	7.05	9.3
	30	265.72	28.08	9.5
	32	1057.77	107.26	9.9

Generation of all subsets of a set of given size

	size(n,k)	sequential time	parallel time	speedup
		(seconds)	(seconds)	(with 16 CPUs)
lexicographic:				
	24,12	1.05	0.23	4.6
	$26,\!13$	4.06	0.62	6.5
	28,14	15.68	2.14	7.3
	$30,\!15$	60.63	8.08	7.5
	32,16	242.23	31.00	7.8
	$34,\!17$	912.48	119.49	7.6
gray code:				
	24,12	0.72	0.17	4.2
	$26,\!13$	2.79	0.37	7.5
	28,14	10.78	1.19	9.1
	$30,\!15$	41.66	4.33	9.6
	32,16	161.67	16.75	9.7
	34,17	628.63	63.99	9.8

Generation of all k-subsets of a set of size n

Experimental Setup

• The *sequential programs* were timed in the following environment:

The compiler used was the GNU C++ compiler. The specific version used is specified below.

g++ --version egcs-2.90.29 980515 (egcs-1.0.3 release)

The programs were compiled and run under Red Hat Linux (version 5.2) running a 2.0.36 kernel (recompiled for SMP support).

```
uname -a
Linux kohinoor 2.0.36 #1 Sat Apr 10 17:46:43 MDT 1999 i586 unknown
```

The machine used is a dual-processor with 166 MHz Pentiums with 64M of RAM.

• The *parallel programs* were implemented using PVM version 3.4.1 and timed on a 8-machine, 16-processor Beowulf cluster connected by two 100 Mbits/s Fast Ethernet. Each machine is identical to the machine on which the sequential timing was performed.

```
kohinoor:for i in 0 1 2 3 4 5 6 7
> do
> echo "On machine ppO$i";rsh ppO$i pvm <<end
> version
> end
> done
On machine pp00
pvm> 3.4.1
On machine pp01
pvm> 3.4.1
On machine pp02
pvm> 3.4.1
On machine pp03
pvm> 3.4.1
On machine pp04
pvm> 3.4.1
On machine pp05
pvm> 3.4.1
On machine pp06
pvm> 3.4.1
On machine pp07
pvm> 3.4.1
```

Further Directions

- The current implementations of ranking and unranking are limited due to the limited precision of integers in C++. We have investigated and found a suitable library for unlimited precision arithmetic that can be used to remove this restriction. The implementations for generating the subsets does not depend on the precision of integers.
- Add more families of combinatorial objects to the library.
- A surprising discovery was that the speedup was not linear as was expected. A further direction would be to investigate how to improve the speedup.